

OS - 2 marks

Q Objectives of OS:

- Make computer system convenient to use in an efficient manner.
- Provides a convenient interface to use the computer system.
- Easy to manage.
- Multitasking.

Q Components of process control blocks:

- Process state
- Process number
- Program counter
- Register
- Memory limits
- Accounting information
- I/O status information.

- (B) A multithreaded system consisting of multiple user-level threads mapped to one kernel thread can't use of the different processors in multiprocessor system. Consequently, there is no performance benefit

associated with this solution. The multithreaded solution could be faster if the multiple user-level threads are mapped to different kernel threads.

Q Benefits of thread:

Responsiveness

Resource sharing

Economy of users resources

Scalability

challenges of thread:

Identifying tasks

Balance

Data splitting

Data dependency

Testing and debugging.

⑤ Execution of bootstrap program:

A bootstrap program is the first code that is executed when the computer system is started. The entire OS depends on the bootstrap program to work correctly as it loads the OS.

⑥ Throughput:

One measure of work is no. of processes that are completed per time unit called throughput.

Turnaround time:

The interval from time of submission of a process to the time of completion is called turnaround time.

CPU utilization:

It can range from 0% to 100%. In real system, it should range from 40% to 90%.

⑦ Busy waiting:

It is a process synchronization technique in which a process (task) waits and constantly checks for a condition to be satisfied before proceeding with its execution.

⑧ primitive operation of semaphore:

wait(s) - Denoted (P) → Proberen means test

signals(s) - Denoted (V) → Verhogen means increment.

⑨ conditions for deadlock:

• mutual exclusion: Atleast one process held in non-shareable state.

• Hold & wait: A process must be holding atleast 1 resource and waiting for additional resource that is currently being held by other process.

• No-preemption: Resource can't be preempted before the process has complete its task.

• circular waits: Set of waiting processes must exist such that P₀ waits for resource held by P₁, P₁ waits for " " P₂, so on.

⑩ Race condition:

Several processes access and manipulate the same data concurrently and outcome of execution depends on the particular order in which the access takes place is called race condition.

⑪ Internal fragmentation

External fragmentation.

⑫ Load time address binding:

If the address is not known at compile time that the process will reside in memory then the compiler must generate relocatable code.

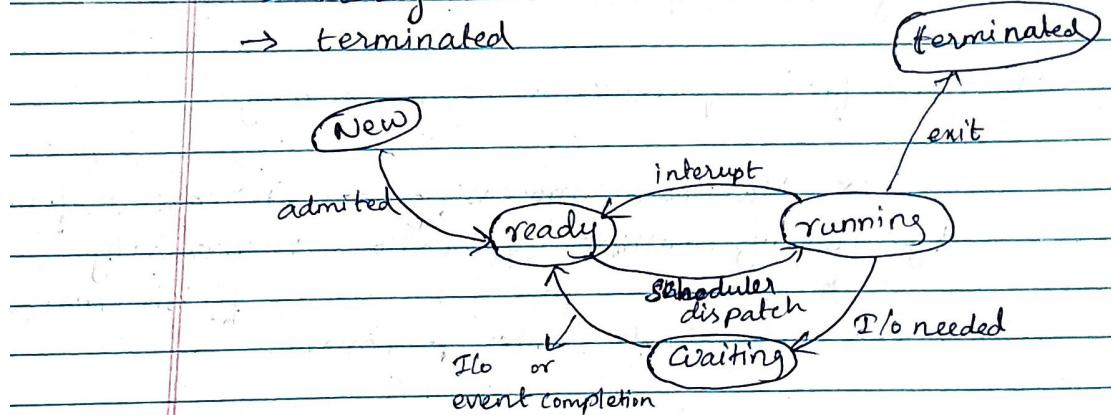
2 Marks

13. Operating System:

An Operating System is a program that manages the computer hardware. It also provides the bases for application programs and act as an intermediate between the computer user and the computer hardware.

14. Process State:

- new
- running
- waiting
- ready
- terminated



15. Dual mode operations:

To ensure the proper executing of the OS we must be able to distinguish between the execution of OS mode and user defined mode.

We need to separate mode of operation. User mode and Kernel mode.

16. Working principle of system boot:

1. Power on Self Test (POST) :

When the computer is turned on, the basic I/O system performs a series of self test to check the system hardware.

2. Boot loader :

It is responsible for loading the operating system kernel into memory.

3. Kernel Initialization:

After boot loader kernel initializes the system hardware and device drivers.

4. User mode - User can log in and start using the computer.

17. Benefits of multicore programming:

* Multicore processing can increase performance by running multiple application concurrently.

* Multicore processors can finish more work than single core processors.

* Performance

* Reliability

* Software interactions

* Multitasking

* Power consumption.

18. Four essential section in process

Synchronization:

- * Entry section
- * Critical Section
- * Exit Section
- * Remainder Section.

19.

TAT

WT

* It refers to the time taken by a process since it enters a ready queue for the process of execution till completion.

It refers to the total time that a process spends while waiting in a ready queue until it gets the CPU.

* Speed of O/P is effective

* Speed of O/P is not effective.

* $TAT = CT - AT$

* $WT = TAT - BT$.

20. Requirements to solve critical Section problem

- * Mutual exclusion
- * Progress,
- * Bounded waiting.

21. Deadlock can be avoided by

21. Deadlock avoidance can be done with Banker's algorithm and Resource allocation graph

 - * Banker's algorithm - when resource have multiple instances, we use banker's algorithm
 - * Resource allocation graph - is the pictorial view of all allocation resources, available resource and its current state

| Logical address | Physical address |
|--|--|
| * Generated by CPU | * Location in memory unit. |
| * User can view the logical address of the program. | * User can never view the physical address of program. |
| * It can be changed. | * It can't be changed. |
| * Also called " virtual " "Virtual address" | * It is also called "real address". |

23. External fragmentation:

 - * External fragmentation occurs whenever a method of dynamic memory allocation happens to allocate some memory and leave a small amount of unusable memory.
 - * The total quantity of the memory available is reduced substantially in case there's too much external fragmentation.